

# 4Gb NAND FLASH HY27UT084G2M



## Document Title 4Gbit (512Mx8bit) NAND Flash Memory

# **Revision History**

Revision No.	I	Draft Date	Remark		
0.0	Initial Draft.	Initial Draft.			
0.1	Delete Double Page Progration     Texts & Figures are deleted     Delete Copy-back Operation     Correct Program Time (tPleter temperature)     temperature temperature     temperature temperature     temperature temperature     temperature temperature     temperatu	Apr. 15. 2005	Preliminary		
0.2	1) Change AC Conditions tab 2) Add tWW parameter ( tW' - Texts & Figures are added tWW is added in AC timing 3) Edit System Interface Usir 4) Correct Program Time (tR)  tPROG Typ.  Before 600us  After 800us  5) Correct Address Cycle May 6) Add Read ID table 7) Correct PKG dimension (TC)  CP  Before 0.050  After 0.100  8) Delete ECC	W = 100ns, not characteristic agree (TE don't control of the POG)  Max.  800us  2000us	s table.	Oct. 12. 2005	Preliminary
0.3	1) Delete Auto Read Page 0 to - The texts & figures are deleter. The definition of PRE is deleter. The definition of PRE is deleter. The definition of PRE is deleter. Of the definition of PRE is deleter. The definition of PRE is deleter. The texts & figures are deleter.	eted. eted. . (Table 9) 512byte) oction		Dec. 23. 2005	Preliminary



# **Revision History**

-Continued-

Revision No.		History					Draft Date	Remark			
	1) Chang	e AC tiı	ming cha	aracteris	stics						
		tWP	tWH	tWC	tALS	tALH	tDH	tR	tAR		
	Before	25	15	50	0	10	10	50	10		
	After	40	20	60	5	15	15	60	20		
		tCLR	tRC	tREA	tREH	tCEA	tRP	tWHR	tRST		
	Before	10	50	30	15	45	25	60	5/10/500		
	After	20	60	45	20	55	40	80	20/50/500		
0.4	2) Chang 3) Chang - Number 4) Chang	e NOP of par	tial prog	ram cyc		e same	page is	s chang	ed to 1.	Jan. 18. 2006	Preliminary
			valid b	valid block number							
	Befo	ore		2008							
	Afte	er		1968							
	1) Chang	e valid	block nu	ımber (	min)						
			valid b	lock nu	mber						
0.5	Befo	ore		1968						Feb. 20. 2006	
	Afte	er		1989							
	2) Delete Preliminary.										
	1) Edit DC and Operating Characteristics (Table 9)										
		Tes	t Conditions (ICC4)			Test Conditions (ICC5)			C5)		
0.6	Before	PR		CE=VIH, E=0V, WP=0V, Vcc			CE=Vcc-0.2, PRE=0V, WP=0V, Vcc			Feb. 28. 2006	
	After			VIH, V, Vcc		CE=Vcc-0.2, WP=0V, Vcc					
	2) Correc	t Read	ID Cycle	e namin	g.						





#### **FEATURES SUMMARY**

#### HIGH DENSITY NAND FLASH MEMORIES

- Cost effective solutions for mass storage applications

#### NAND INTERFACE

- x8 bus width
- Multiplexed address/ Data
- Pinout compatibility for all densities

#### **SUPPLY VOLTAGE**

- 3.3V device : VCC = 2.7V to 3.6 V : HY27UT084G2M

#### **MEMORY CELL ARRAY**

- (2K + 64) bytes x 128 pages x 2048 blocks

#### **PAGE SIZE**

x8 device : (2048+64 spare) bytes: HY27UT084G2M

#### **BLOCK SIZE**

- x8 device : (256K+8K) bytes

#### **PAGE READ / PROGRAM**

Random access: 60us (Max)Sequentiall access: 60ns (Min)Page program time: 800us (Typ)

#### **FAST BLOCK ERASE**

- Block erase time: 2ms (Typ)

#### STATUS REGISTER

#### **ELECTRONIC SIGNATURE**

- 1st cycle: Manufacturer Code
- 2nd cycle: Device Code
- 3rd cycle: Internal chip number, Cell Type, Number of

Simultaneously Programmed Pages.

- 4th cycle: Page size, Block size, Organization, Spare

size

#### CHIP ENABLE DON'T CARE

-Simple interface with microcontroller

#### HARDWARE DATA PROTECTION

- Program/Erase locked during Power transitions.

#### **SERIAL NUMBER OPTION**

#### **DATA RETENTION**

- 10000 Program / Erase cycles (with 4bit/528byte ECC)
- 10 Years Data Retention

#### **PACKAGE**

- HY27UT084G2M-T(P)
  - : 48-pin TSOP1(12 x 20 x 1.2 mm)
    - HY27UT084G2M-T (Lead)
  - HY27UT084G2M-TP (Lead Free)



#### 1. SUMMARY DESCRIPTION

HY27UT084G2M is a 512Mx8bit with spare 16Mx8bit capacity. The device is offered in 3.3V Vcc Power supply. Its MLC NAND cell provides the most cost-effective solution for the solid state mass storage market.

The memory is divided into blocks that can be erased independently so it is possible to preserve valid data while old

The device contains 2048 blocks, composed by 128 pages consisting in two NAND structures of 32 series connected Flash cells. All cell hold two bit. Even pages are stored in the LSB of the cells and odd pages are stored in the MSB of the cells.

A program operation allows to write the 2112-byte page in typical 800us and an erase operation can be performed in typical 2ms on a 256K-byte block.

Data in the page can be read out at 60ns cycle time per byte. The I/O pins serve as the ports for address and data input/output as well as command input. This interface allows a reduced pin count and easy migration towards different densities, without any rearrangement of footprint.

Commands, Data and Addresses are synchronously introduced using  $\overline{\text{CE}}$ ,  $\overline{\text{WE}}$ , ALE and CLE input pin. The on-chip Program/Erase Controller automates all program and erase functions including pulse repetition, where required, and internal verification and margining of data.

The modify operations can be locked using the WP input pin.

The output pin R/B (open drain buffer) signals the status of the device during each operation. In a system with multiple memories the R/B pins can be connected all together to provide a global status signal.

Even the write-intensive systems can take advantage of the HY27UT084G2M's extended reliability of 10K (with 4bit/ 528 Byte ECC) Program/Erase cycles.

The chip is offered with the CE don't care function. This option allows the direct download of the code from the NAND Flash memory device by a microcontroller, since the CE transitions do not stop the read operation.

The HY27UT084G2M is available in 48 - TSOP1 12 x 20 mm package.

#### 1.1 Product List

PART NUMBER	ORIZATION	VCC RANGE	PACKAGE
HY27UT084G2M	x8	2.7~3.6 Volt	TSOP48



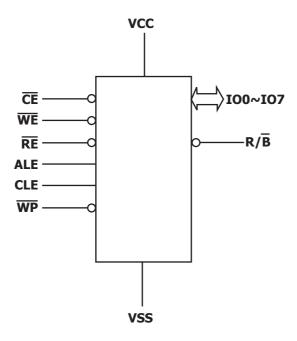


Figure1: Logic Diagram

IO15 - IO8	Data Input / Outputs
107 - 100	Data Input / Outputs
CLE	Command latch enable
ALE	Address latch enable
CE	Chip Enable
RE	Read Enable
WE	Write Enable
WP	Write Protect
R/B	Ready / Busy
Vcc	Power Supply
Vss	Ground
NC	No Connection

Table 1: Signal Names



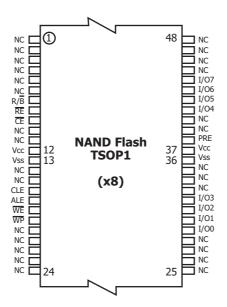


Figure 2. 48TSOP1 Contactions, x8 Device



#### 1.2 PIN DESCRIPTION

Pin Name	Description
100-107	DATA INPUTS/OUTPUTS  The IO pins allow to input command, address and data and to output <u>data</u> during read / program operations. The inputs are latched on the rising edge of Write Enable (WE). The I/O buffer float to High-Z when the device is deselected or the outputs are disabled.
CLE	COMMAND LATCH ENABLE This input activates the latching of the IO inputs inside the Command Register on the Rising edge of Write Enable (WE).
ALE	ADDRESS LATCH ENABLE This input activates the latching of the IO inputs inside the Address Register on the Rising edge of Write Enable (WE).
CE	CHIP ENABLE  This input controls the selection of the device. When the device is busy $\overline{\text{CE}}$ low does not deselect the memory.
WE	WRITE ENABLE This input acts as clock to latch Command, Address and Data. The IO inputs are latched on the rise edge of WE.
RE	<b>READ ENABLE</b> The RE input is the serial data-out control, and when active drives the data onto the I/O bus. Data is valid tREA after the falling edge of RE which also increments the internal column address counter by one.
WP	WRITE PROTECT The WP pin, when Low, provides an Hardware protection against undesired modify (program / erase) operations.
R/B	READY BUSY The Ready/Busy output is an Open Drain pin that signals the state of the memory.
VCC	SUPPLY VOLTAGE The VCC supplies the power for all the operations (Read, Write, Erase).
VSS	GROUND
NC	NO CONNECTION

## **Table 2: Pin Description**

#### NOTE:

1. A 0.1uF capacitor should be connected between the VCC Supply Voltage pin and the VSS Ground pin to decouple the current surges from the power supply. The PCB track widths must be sufficient to carry the currents required during program and erase operations.



	100	I01	102	103	104	105	106	107
1st Cycle	A0	A1	A2	А3	A4	<b>A</b> 5	A6	A7
2nd Cycle	A8	A9	A10	A11	L <sup>(1)</sup>	L <sup>(1)</sup>	L <sup>(1)</sup>	L <sup>(1)</sup>
3rd Cycle	A12	A13	A14	A15	A16	A17	A18	A19
4th Cycle	A20	A21	A22	A23	A24	A25	A26	A27
5th Cycle	A28	A29	L <sup>(1)</sup>					

Table 3: Address Cycle Map(x8)

#### NOTE:

1. L must be set to Low.

FUNCTION	1st CYCLE	2nd CYCLE	3rd CYCLE	Acceptable command during busy
READ 1	00h	30h	-	
READ ID	90h	-	-	
RESET	FFh	-	-	Yes
PAGE PROGRAM (start)	80h	10h	-	
BLOCK ERASE	60h	D0h	-	
READ STATUS REGISTER	70h	-	-	Yes
RANDOM DATA INPUT	85h	-	-	
RANDOM DATA OUTPUT	05h	E0h	-	

**Table 4: Command Set** 



CLE	ALE	CE	WE	RE	WP		MODE
Н	L	L	Rising	Н	Х	Read Mode	Command Input
L	Н	L	Rising	Н	Х	Read Mode	Address Input(5 cycles)
Н	L	L	Rising	Н	Н	Write Mode	Command Input
L	Н	L	Rising	Н	Н	write wode	Address Input(5 cycles)
L	L	L	Rising	Н	Н	Data Input	
L	L	L <sup>(1)</sup>	Н	Falling	Х	Sequential Read and Data Output	
L	L	L	Н	Н	Х	During Read (	(Busy)
Х	Х	Х	Х	Х	Н	During Progra	m (Busy)
Х	Х	Х	Х	Х	Н	During Erase (Busy)	
Х	Х	Х	Х	Х	L	Write Protect	
Х	Х	Н	Х	Х	0V/Vcc	Stand By	

**Table 5: Mode Selection** 

#### NOTE:

1. With the  $\overline{\text{CE}}$  don't care option  $\overline{\text{CE}}$  high during latency time does not stop the read operation



#### 2. BUS OPERATION

There are six standard bus operations that control the device. These are Command Input, Address Input, Data Input, Data Output, Write Protect, and Standby.

Typically glitches less than 5 ns on Chip Enable, Write Enable and Read Enable are ignored by the memory and do not affect bus operations.

#### 2.1 Command Input.

Command Input bus operation is used to give a command to the memory device. Commands are accepted with Chip Enable low, Command Latch Enable High, Address Latch Enable low and Read Enable High and latched on the rising edge of Write Enable. Moreover for commands that start a modify operation (write/erase) the Write Protect pin must be high. See figure 4 and table 14 for details of the timings requirements. Command codes are always applied on IO7:0, disregarding the bus configuration (X8).

#### 2.2 Address Input.

Address Input bus operation allows the insertion of the memory address. To insert the 30 addresses is needed to access the 4Gbit, 5 clock cycles are needed. Addresses are accepted with Chip Enable low, Address Latch Enable High, Command Latch Enable low and Read Enable High and latched on the rising edge of Write Enable. Moreover for commands that start a modify operation (write/erase) the Write Protect pin must be high. See figure 5 and table 15 for details of the timings requirements. Addresses are always applied on IO7:0, disregarding the bus configuration (X8).

#### 2.3 Data Input.

Data Input bus operation allows to feed to the device the data to be programmed. The data insertion is serially and timed by the Write Enable cycles. Data are accepted only with Chip Enable low, Address Latch Enable low, Command Latch Enable low, Read Enable High, and Write Protect High and latched on the rising edge of Write Enable. See figure 6 and table 14 for details of the timings requirements.

#### 2.4 Data Output.

Data Output bus operation allows to read data from the memory array and to check the status register content, the ID data. Data can be serially shifted out toggling the Read Enable pin with Chip Enable low, Write Enable High, Address Latch Enable low, and Command Latch Enable low. See figures 7,9,10 and table 14 for details of the timings requirements.

#### 2.5 Write Protect.

Hardware Write Protection is activated when the Write Protect pin is low. In this condition modify operation do not start and the content of the memory is not altered. Write Protect pin is not latched by Write Enable to ensure the protection even during the power up.

#### 2.6 Standby.

In Standby the device is deselected, outputs are disabled and Power Consumption reduced.



#### 3. DEVICE OPERATION

#### 3.1 Page Read.

Upon initial device power up, the device defaults to Read mode. This operation is also initiated by writing 00h and 30h to the command register along with five address cycles. In two consecutive read operations, the second one doesn't' need 00h command, which five address cycles and 30h command initiates that operation. Two types of operations are available: random read, serial page read. The random read mode is enabled when the page address is changed. The 2112 bytes (X8 device) of data within the selected page are transferred to the data registers in less than 60us(tR). The system controller may detect the completion of this data transfer 60us(tR) by analyzing the output of R/B pin. Once the data in a page is loaded into the data registers, they may be read out in 60ns cycle time by sequentially pulsing  $\overline{\text{RE}}$ . The repetitive high to low transitions of the  $\overline{\text{RE}}$  clock make the device output the data starting from the selected column address up to the last column address.

The device may output random data in a page instead of the consecutive sequential data by writing random data output command.

The column address of next data, which is going to be out, may be changed to the address which follows random data output command.

Random data output can be operated multiple times regardless of how many times it is done in a page.

#### 3.2 Page Program.

The device is programmed by page The number of consecutive partial page programming operation within the same page without an intervening erase operation must not exceed 1 time.

The addressing should be done in sequential order in a block. A page program cycle consists of a serial data loading period in which up to 2112bytes (X8 device) of data may be loaded into the data register, followed by a non-volatile programming period where the loaded data is programmed into the appropriate cell.

The serial data loading period begins by inputting the Serial Data Input command (80h), followed by the five cycle address inputs and then serial data. The words other than those to be programmed do not need to be loaded. The device supports random data input in a page. The column address of next data, which will be entered, may be changed to the address which follows random data input command (85h). Random data input may be operated multiple times regardless of how many times it is done in a page.

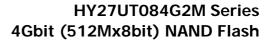
The Page Program confirm command (10h) initiates the programming process. Writing 10h alone without previously entering the serial data will not initiate the programming process.

The internal write state controller automatically executes the algorithms and timings necessary for program and verify, thereby freeing the system controller for other tasks. Once the program process starts, the Read Status Register command may be entered to read the status register.

The system controller can detect the completion of a program cycle by monitoring the  $R/\overline{B}$  output, or the Status bit (I/O 6) of the Status Register. Only the Read Status command and Reset command are valid while programming is in progress. When the Page Program is complete, the Write Status Bit (I/O 0) may be checked.

The internal write verify detects only errors for "1"s that are not successfully programmed to "0"s.

The command register remains in Read Status command mode until another valid command is written to the command register. Figure 12 details the sequence.





#### 3.3 Block Erase.

The Erase operation is done on a block basis. Block address loading is accomplished in two cycles initiated by an Erase Setup command (60h). Only address A19 to A29 (X8) is valid while A12 to A18 (X8) is ignored. The Erase Confirm command (D0h) following the block address loading initiates the internal erasing process. This two-step sequence of setup followed by execution command ensures that memory contents are not accidentally erased due to external noise conditions.

At the rising edge of  $\overline{\text{WE}}$  after the erase confirm command input, the internal write controller handles erase and erase-verify.

Once the erase process starts, the Read Status Register command may be entered to read the status register. The system controller can detect the completion of an erase by monitoring the  $R/\overline{B}$  output, or the Status bit (I/O 6) of the Status Register. Only the Read Status command and Reset command are valid while erasing is in progress. When the erase operation is completed, the Write Status Bit (I/O 0) may be checked.

Figure 14 details the sequence.

#### 3.4 Read Status Register.

The device contains a Status Register which may be read to find out whether read, program or erase operation is completed, and whether the program or erase operation is completed successfully. After writing 70h command to the command register, a read cycle outputs the content of the Status Register to the I/O pins on the falling edge of  $\overline{\text{CE}}$  or  $\overline{\text{RE}}$ , whichever occurs last. This two line control allows the system to poll the progress of each device in multiple memory connections even when  $R/\overline{B}$  pins are common-wired.  $\overline{\text{RE}}$  or  $\overline{\text{CE}}$  does not need to be toggled for updated status. Refer to table 15 for specific Status Register definitions. The command register remains in Status Read mode until further commands are issued to it. Therefore, if the status register is read during a random read cycle, the read command (00h) should be given before starting read cycles.

#### 3.5 Read ID.

The device contains a product identification mode, initiated by writing 90h to the command register, followed by an address input of 00h. Four read cycles sequentially output the manufacturer code (ADh), and the device code and 3rd cycle ID, 4th cycle ID respectively. The command register remains in Read ID mode until further commands are issued to it. Figure 15 shows the operation sequence, while tables 17 explain the byte meaning.

#### 3.6 Reset.

The device offers a reset feature, executed by writing FFh to the command register. When the device is in Busy state during random read, program or erase mode, the reset operation will abort these operations. The contents of memory cells being altered are no longer valid, as the data will be partially programmed or erased. The command register is cleared to wait for the next command, and the Status Register is cleared to value E0h when  $\overline{\text{WP}}$  is high. Refer to table 15 for device status after reset operation. If the device is already in reset state a new reset command will not be accepted by the command register. The R/ $\overline{\text{B}}$  pin goes low for tRST after the Reset command is written. Refer to Figure 18.



#### 4. OTHER FEATURES

#### 4.1 Data Protection & Power On/Off Sequence

The device is designed to offer protection from any involuntary program/erase during power-transitions. An internal voltage detector disables all functions whenever Vcc is below about 2V(3.3V device). WP pin provides hardware protection and is recommended to be kept at VIL during power-up and power-down. A recovery time of minimum 10us is required before internal circuit gets ready for any command sequences as shown in Figure 19. The two-step command sequence for program/erase provides additional software protection.

If the power is dropped during the ready read/write/erase operation, Power protection function may not guaranteed the data. Power protection function is only available during the power on/off sequence.

#### 4.2 Ready/Busy.

The device has a Ready/Busy output that provides method of indicating the completion of a page program, erase, random read completion. The  $R/\overline{B}$  pin is normally high and goes to low when the device is busy (after a reset, read, program, erase operation). It returns to high when the internal controller has finished the operation. The pin is an opendrain driver thereby allowing two or more  $R/\overline{B}$  outputs to be Or-tied. Because pull-up resistor value is related to  $tr(R/\overline{B})$  and current drain during busy (Ibusy), an appropriate value can be obtained with the following reference chart (Fig 20). Its value can be determined by the following guidance.



Parameter	Symbol	Min	Тур	Max	Unit
Valid Block Number	NvB	1989		2048	Blocks

**Table 6: Valid Blocks Number** 

#### NOTE:

1. The 1st block is guaranteed to be a valid block at the time of shipment.

Symbol	Parame	Parameter		
	Ambient Operating Temperature (Te	mperature Range Option 1)	0 to 70	$^{\circ}$
Та	Ambient Operating Temperature (Te	mperature Range Option 2)	-25 to 85	$^{\circ}$
	Ambient Operating Temperature (Te	-40 to 85	$^{\circ}$ C	
TBIAS	Temperature Under Bias	-50 to 125	$^{\circ}$	
Tstg	Storage Temperature		-65 to 150	$^{\circ}$
Vio <sup>(2)</sup>	Input or Output Voltage 3.3V Device		-0.6 to 4.6	V
Vcc	Supply Voltage	-0.6 to 4.6	V	

**Table 7: Absolute maximum ratings** 

#### NOTE:

- 1. Except for the rating "Operating Temperature Range", stresses above those listed in the Table "Absolute Maximum Ratings" may cause permanent damage to the device. These are stress ratings only and operation of the device at these or any other conditions above those indicated in the Operating sections of this specification is not implied. Exposure to Absolute Maximum Rating conditions for extended periods may affect device reliability. Refer also to the Hynix SURE Program and other relevant quality documents.
- 2. Minimum Voltage may undershoot to -2V during transition and for less than 20ns during transitions.



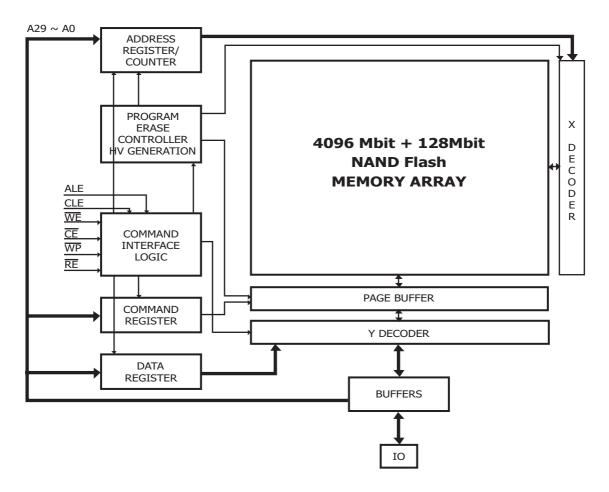


Figure 3: Block Diagram



Parameter		Compleal	Took Conditions		11		
Рага	ameter	Symbol	Test Conditions	Min	Тур	Max	Unit
Operating	Sequential Read	Icc1	trc=60ns CE=VIL, IOUT=0mA	-	15	30	mA
Current	Program	ICC2	-	-	15	30	mA
	Erase	Icc3	-	-	15	30	mA
Stand-by Curre	nt (TTL)	ICC4	<u>CE</u> =VIH, <u>WP</u> =0V, Vcc	-		1	mA
Stand-by Curre	nt (CMOS)	ICC5	CE=Vcc-0.2, WP=0V, Vcc	-	10	50	uA
Input Leakage	Current	ILI	VIN=0 to Vcc (max)	-	-	± 10	uA
Output Leakage	Current	ILO	Vout =0 to Vcc (max)	-	-	± 10	uA
Input High Volta	age	ViH	-	0.8xVcc	-	Vcc+0.3	V
Input Low Volta	ige	VIL	-	-0.3	-	0.2xVcc	V
Output High Vo	Itage Level	Vон	Ioн=-400uA	2.4	-	-	V
Output Low Vol	tage Level	Vol	IoL=2.1mA	-	-	0.4	V
Output Low Cur	rent (R/B)	Io <u>L</u> (R/B)	Vol=0.6V	8	10	-	mA

**Table 8: DC and Operating Characteristics** 

Parameter	Value	
Parameter	3.3Volt	
Input Pulse Levels	0.4V to 2.4V	
Input Rise and Fall Times	5ns	
Input and Output Timing Levels	1.5V	
Output Load (2.7V - 3.3V)	1 TTL GATE and CL=50pF	
Output Load (3.0V - 3.6V)	1 TTL GATE and CL=100pF	

**Table 9: AC Conditions** 



Item	Symbol	Test Condition	Min	Max	Unit
Input / Output Capacitance	C1/0	VIL=0V	-	10	pF
Input Capacitance	CIN	VIN=0V	-	10	pF

Table 10: Pin Capacitance (TA=25C, F=1.0MHz)

Parameter	Symbol	Min	Тур	Max	Unit
Program Time	tprog	-	800	2000	us
Number of partial Program Cycles in the same page	NOP	-	-	1	Cycles
Block Erase Time	tBERS	-	2	3	ms

**Table 11: Program / Erase Characteristics** 



# HY27UT084G2M Series 4Gbit (512Mx8bit) NAND Flash

Paramatan.	Complete I	3.3	Volt	11
Parameter	Symbol	Min	Max	Unit
CLE Setup time	tcls	0		ns
CLE Hold time	tсьн	10		ns
CE setup time	tcs	0		ns
CE hold time	tсн	10		ns
WE pulse width	twp	40		ns
ALE setup time	tals	5		ns
ALE hold time	talh	15		ns
Data setup time	tos	20		ns
Data hold time	tон	15		ns
Write Cycle time	twc	60		ns
WE High hold time	twн	20		ns
Data Transfer from Cell to register	tr		60	us
ALE to RE Delay	tar	20		ns
CLE to RE Delay	tclr	20		ns
Ready to RE Low	trr	20		ns
RE Pulse Width	trp	40		ns
WE High to Busy	twB		100	ns
Read Cycle Time	trc	60		ns
RE Access Time	trea		45	ns
RE High to Output High Z	trhz		30	ns
CE High to Output High Z	tcHz		20	ns
RE or CE High to Output Hold	toн	10		ns
RE High Hold Time	treh	20		ns
Output High Z to RE low	tır	0		ns
CE Access Time	tcea		55	ns
WE High to RE low	twhr	80		ns
Device Resetting Time (Read / Program / Erase)	trst		20/50/500 <sup>(1)</sup>	us
Write Protection time	tww <sup>(2)</sup>	100		ns

**Table 12: AC Timing Characteristics** 

#### NOTE:

- 1. If Reset Command (FFh) is written at Ready state, the device goes into Busy for maximum 5us
- 2. Program / Erase Enable Operation : WP high to WE High.

  Program / Erase Disable Operation : WP Low to WE High.



10	Pagae Program	Block Erase	Read	CODING
0	Pass / Fail	Pass / Fail	NA	Pass: '0' Fail: '1'
1	NA	NA	NA	-
2	NA	NA	NA	-
3	NA	NA	NA	-
4	NA	NA	NA	-
5	Ready/Busy	Ready/Busy	Ready/Busy	Active: '0' Idle: '1'
6	Ready/Busy	Ready/Busy	Ready/Busy	Busy: '0' Ready': '1'
7	Write Protect	Write Protect	Write Protect	Protected: '0' Not Protected: '1'

**Table 13: Status Register Coding** 

DEVICE IDENTIFIER CYCLE	DESCRIPTION
1st	Manufacturer Code
2nd	Device Identifier
3rd	Internal chip number, cell Type, Number of Simultaneously Programmed pages.
4th	Page Size, Block Size, Spare Size, Organization

**Table 14: Device Identifier Coding** 

Part Number	Voltage	Bus Width	1st cycle (Manufacture Code)	2nd cycle (Device Code)	3rd cycle	4th cycle
HY27UT084G2M	3.3V	х8	ADh	DCh	84h	25h

Table 15: Read ID Data Table



	Description	107	106	105 104	103 102	I01 I00
Die / Package	1 2 4 Reserved					0 0 0 1 1 0 1 1
String Type	Single Level 2x Multi-level Reserved Reserved				0 0 0 1 1 0 1 1	
Number of Simultaneously Programmed Pages	1 2 3 4			0 0 0 1 1 0 1 1		
Interleave Program Between multiple chips	Not Support Support		0 1			
Write Cache	Not Support Support	0 1				

**Table 16: 3rd Byte of Device Idendifier Description** 

	Description	107	106	105-4	103	102	IO1-0
Page Size (Without Spare Area)	1K 2K Reserved Reserved						0 0 0 1 1 0 1 1
Spare Area Size (Byte / 512Byte)	8 16					0 1	
Serial Access Time	50ns / 30ns 25ns Reserved Reserved	0 0 1 1			0 1 0 1		
Block Size (Without Spare Area)	64K 128K 256K Reserved			0 0 0 1 1 0 1 1			
Organization	X8 X16		0 1				

**Table 17: 4th Byte of Device Identifier Description** 



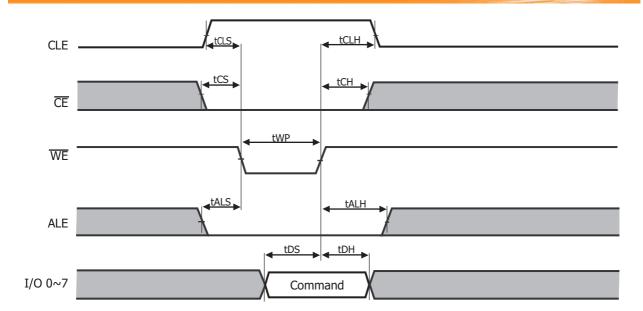


Figure 4: Command Latch Cycle



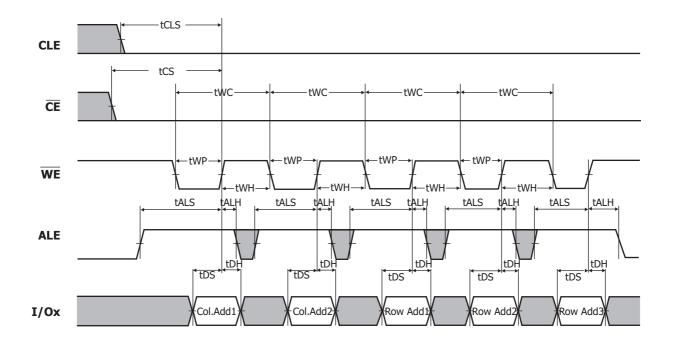


Figure 5: Address Latch Cycle

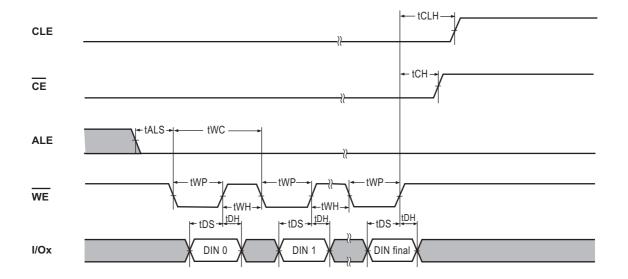
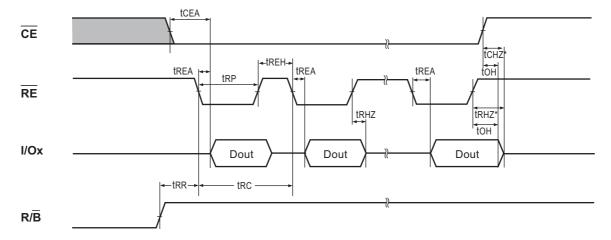


Figure 6. Input Data Latch Cycle





NOTES : Transition is measured ±200mV from steady state voltage with load. This parameter is sampled and not 100% tested.

Figure 7: Sequential Out Cycle after Read (CLE=L, WE=H, ALE=L)

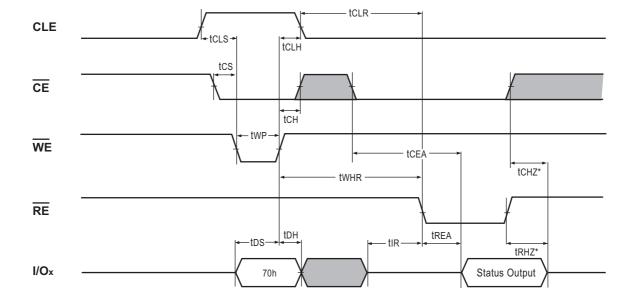


Figure 8: Status Read Cycle



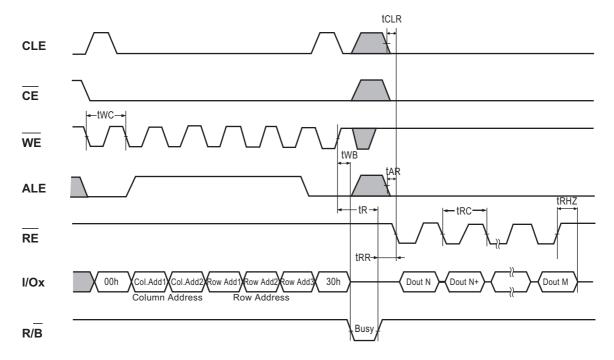


Figure 9: Read1 Operation (Read One Page)

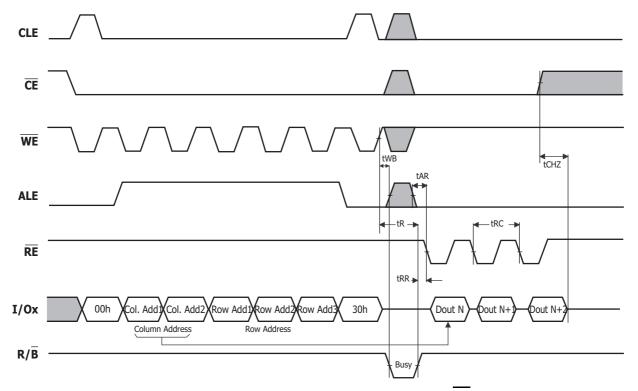


Figure 10: Read1 Operation intercepted by CE

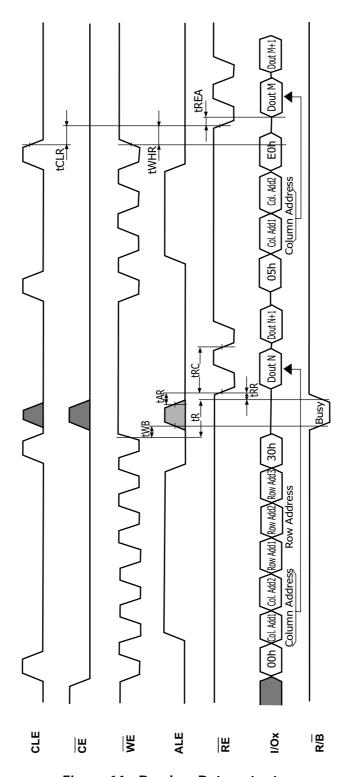
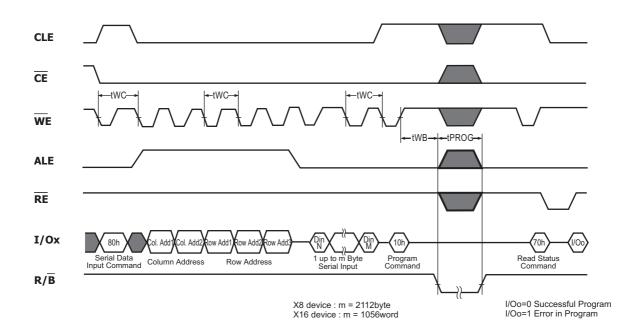


Figure 11 : Random Data output





**Figure 12: Page Program Operation** 

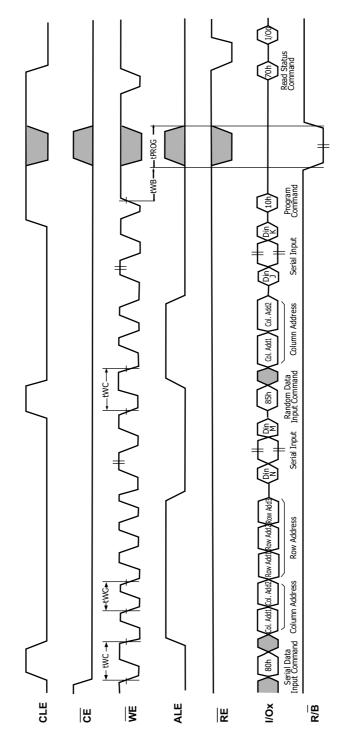


Figure 13: Random Data In



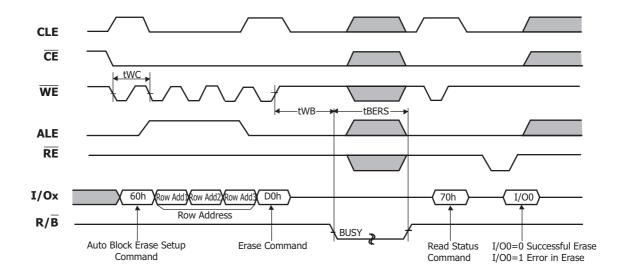


Figure 14: Block Erase Operation (Erase One Block)

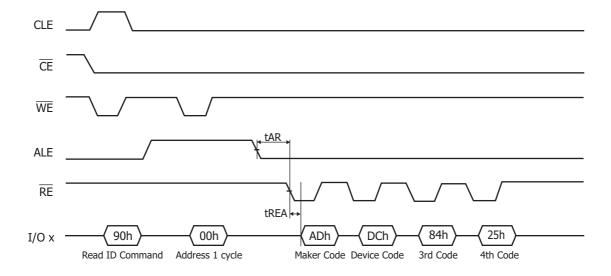


Figure 15: Read ID Operation



#### System Interface Using CE don't care

To simplify system interface,  $\overline{\text{CE}}$  may be deasserted during data loading or sequential data-reading as shown below. So, it is possible to connect NAND Flash to a microporcessor. The only function that was removed from standard NAND Flash to make  $\overline{\text{CE}}$  don't care read operation was disabling of the automatic sequential read function.

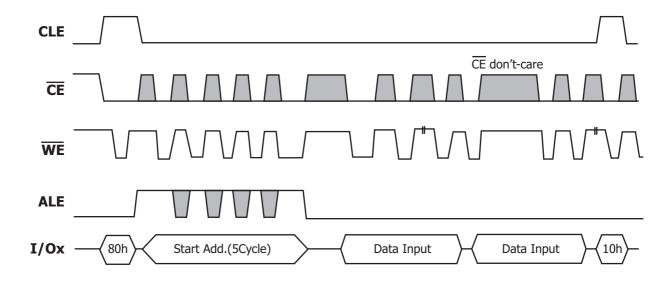


Figure 16: Program Operation with CE don't-care.

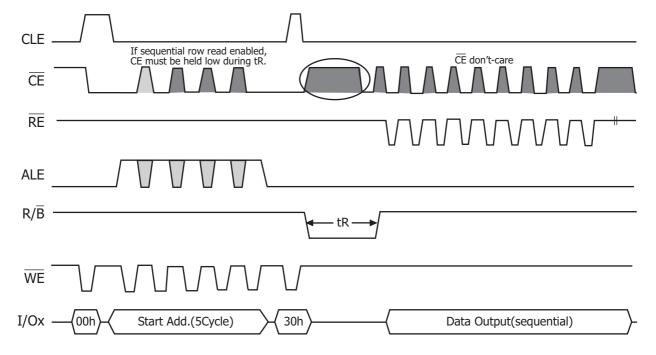


Figure 17: Read Operation with CE don't-care.



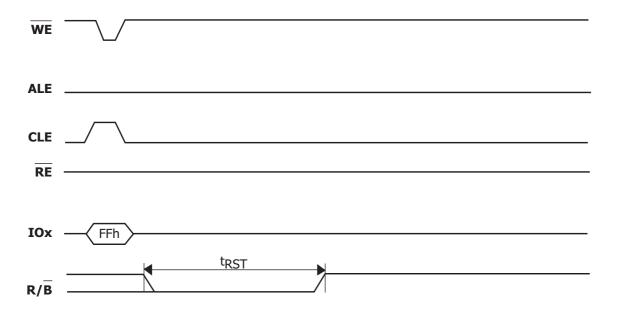


Figure 18: Reset Operation



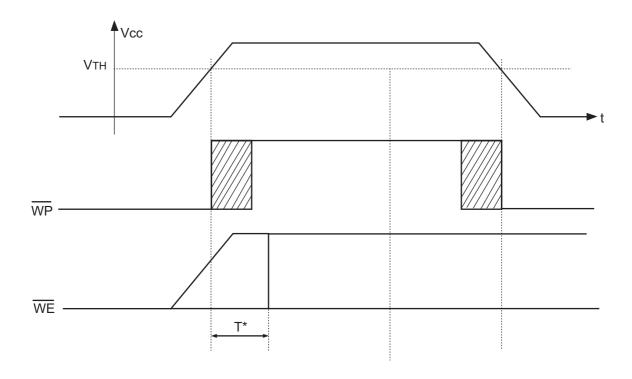


Figure 19: Power On and Data Protection Timing

VTH = 2.5 Volt for 3.3 Volt Supply devices



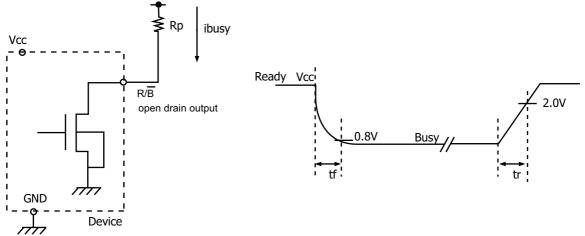
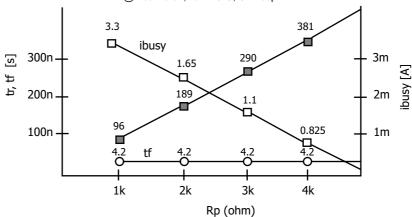


Fig. Rp vs tr, tf & Rp vs ibusy

@ Vcc = 3.3V, Ta = 25°C, CL=100pF



### Rp value guidence

$$Rp (min) = \frac{Vcc (Max.) - Vol (Max.)}{Iol + \Sigma IL} = \frac{3.2V}{8mA + \Sigma IL}$$

where IL is the sum of the input currents of all devices tied to the  $R/\overline{B}$  pin.

Rp(max) is determined by maximum permissible limit of tr

Figure 20: Ready/Busy Pin electrical specifications



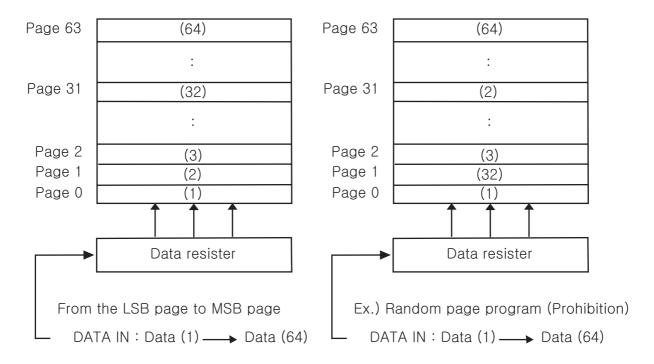


Figure 21: page programming within a block



#### **Bad Block Management**

Devices with Bad Blocks have the same quality level and the same AC and DC characteristics as devices where all the blocks are valid. A Bad Block does not affect the performance of valid blocks because it is isolated from the bit line and common source line by a select transistor. The devices are supplied with all the locations inside valid blocks erased(FFh). The Bad Block Information is written prior to shipping. Any block where the 1st Byte in the spare area of the Last or (Last-2)th page (if the last page is Bad) does not contain FFh is a Bad Block. The Bad Block Information must be read before any erase is attempted as the Bad Block Information may be erased. For the system to be able to recognize the Bad Blocks based on the original information it is recommended to create a Bad Block table following the flowchart shown in Figure 22. The 1st block, which is placed on 00h block address is guaranteed to be a valid block.

#### **Bad Replacement**

Over the lifetime of the device additional Bad Blocks may develop. In this case the block has to be replaced by copying the data to a valid block. These additional Bad Blocks can be identified as attempts to program or erase them will give errors in the Status Register.

As the failure of a page program operation does not affect the data in other pages in the same block, the block can be replaced by re-programming the current data and copying the rest of the replaced block to an available valid block. Refer to Table 18 for the recommended procedure to follow if an error occurs during an operation.

Operation	Recommended Procedure
Erase	Block Replacement
Program	Block Replacement or ECC (with 4bit/528byte)
Read	ECC (with 4bit/528byte)

Table 18: Block Failure

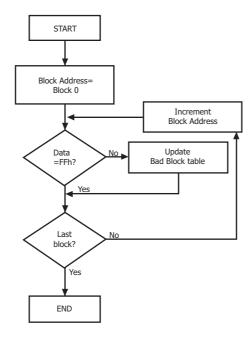


Figure 22: Bad Block Management Flowchart



## **Write Protect Operation**

The Erase and Program Operations are automatically reset when  $\overline{WP}$  goes Low (tWW = 100ns, min). The operations are enabled and disabled as follows (Figure 23~26)

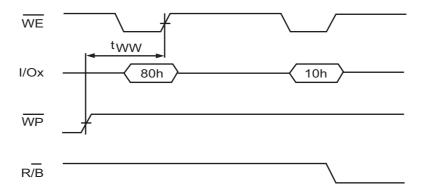


Figure 23: Enable Programming

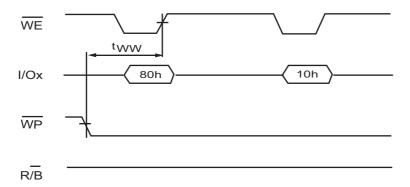


Figure 24: Disable Programming



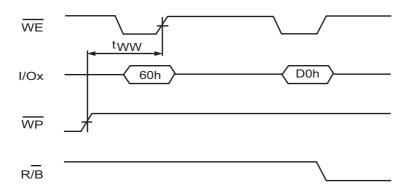


Figure 25: Enable Erasing

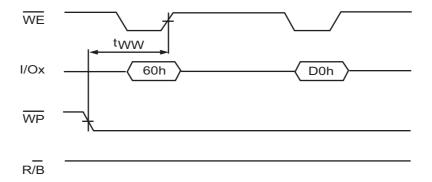


Figure 26: Disable Erasing



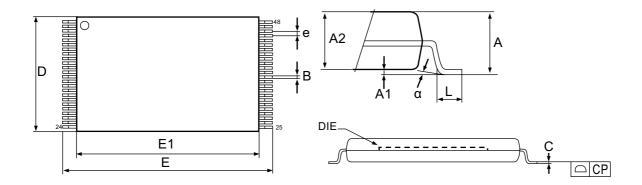


Figure 27. 48-TSOP1 - 48-lead Plastic Thin Small Outline, 12 x 20mm, Package Outline

Symbol		millimeters	
Symbol	Min	Тур	Max
A			1.200
A1	0.050		0.150
A2	0.980		1.030
В	0.170		0.250
С	0.100		0.200
СР			0.100
D	11.910	12.000	12.120
E	19.900	20.000	20.100
E1	18.300	18.400	18.500
е		0.500	
L	0.500		0.680
alpha	0		5

Table 19: 48-TSOP1 - 48-lead Plastic Thin Small Outline, 12 x 20mm, Package Mechanical Data



## **MARKING INFORMATION - TSOP1**

Packag	Marking Example					
TSOP1						
- hynix	: Hynix Symbol					
- KOR	: Origin Country					
- HY27xTxx4Gx HY: Hynix 27: NAND Flas						
x: Power Supp						
T: Classification						
xx: Bit Organiz	ration : 08(x8), 16(x16)					
4G: Density	: 4Gbit					
x: Mode	: 2(1nCE & 1R/nB; Sequential Row Read Disable)					
M: Version	: 1st Generation					
x: Package Typ	De : T(48-TSOP1)					
x: Package Ma	terial : Blank(Normal), P(Lead Free)					
x: Operating T	emperature : C(0℃~70℃), E(-25℃~85℃)					
	M(-30℃~85℃), I(-40℃~85℃)					
x: Bad Block	: B(Included Bad Block), S(1~5 Bad Block),					
	P(All Good Block)					
	year 2005, 06= year 2006)					
	ek (ex: 12= work week 12)					
- xx: Process Co	ae					
Note - Capital Letter	: Fixed Item					
- Small Letter	: Non-fixed Item					
- Jillali Lettel	. NOII-IIAGU IIGIII					